

GAUSSIAN ELIMINATION

$$Ax = b$$

$\left. \begin{array}{l} \text{ } \\ \text{ } \end{array} \right\} \begin{array}{l} n \times n \\ \text{(given)} \end{array}$
 $\left. \begin{array}{l} \text{ } \\ \text{ } \end{array} \right\} \begin{array}{l} \text{in } \mathbb{C}^n \\ \text{(unknown)} \end{array}$
 $\left. \begin{array}{l} \text{ } \\ \text{ } \end{array} \right\} \begin{array}{l} \text{in } \mathbb{C}^n \\ \text{(given)} \end{array}$

where A is nonsingular.

STRATEGY

- (1) Convert $Ax = b$ into an equivalent upper triangular system $Ux = \hat{b}$
- (2) Solve $Ux = \hat{b}$ by back substitution.

DEFN

The systems

$$Ax = b \quad \text{and} \quad \hat{A}x = \hat{b}$$

are called equivalent if they have exactly the same solution set.

Augmented Matrix: $[A | b]$ (for the system $Ax = b$)

Row Reduction

$$[A | b] \xrightarrow{\text{apply row operations repeatedly}} [U | \hat{b}]$$

upper triangular

Row Operations

(1) Row Replace

Add a multiple of a row to another.

$$\left[\begin{array}{cc|c} 2 & 1 & -2 \\ -1 & 3 & 1 \end{array} \right] \xrightarrow{r_2 := r_2 + 2r_1} \left[\begin{array}{cc|c} 2 & 1 & -2 \\ 3 & 5 & -3 \end{array} \right]$$

(2) Row Scale

Multiply a row by a scalar.

$$\left[\begin{array}{cc|c} 2 & 1 & -2 \\ -1 & 3 & 1 \end{array} \right] \xrightarrow{r_2 := -3r_2} \left[\begin{array}{cc|c} 2 & 1 & -2 \\ 3 & -9 & -3 \end{array} \right]$$

~~$r_2 := r_2 + 3r_1$~~

(3) Row Interchange

Swap two rows.

$$\left[\begin{array}{cc|c} 2 & 1 & -2 \\ -1 & 3 & 1 \end{array} \right] \xrightarrow{r_1 \leftrightarrow r_2} \left[\begin{array}{cc|c} -1 & 3 & 1 \\ 2 & 1 & -2 \end{array} \right]$$

BASIC RESULT

Suppose

$$[A \mid b] \xrightarrow{\substack{\text{apply one} \\ \text{of the row operations}}} [\hat{A} \mid \hat{b}]$$

Then $Ax = b$ and $\hat{A}x = \hat{b}$ are equivalent.

$$[A \mid b] \rightsquigarrow [U \mid \hat{b}]$$

apply row operations repeatedly

$$\implies Ax = b \text{ and } Ux = \hat{b}$$

are equivalent systems

ALGORITHM (Gaussian Elimination)

- * Proceed column by column starting with first column moving to the right.
- * Processing i th column
set every entry on the i th column below a_{ii} to zero by applying row replacement operations.

CASE $n=3$

$$\left[\begin{array}{ccc|c} x & x & x & x \\ x & x & x & x \\ x & x & x & x \end{array} \right] \xrightarrow{i=1} \left[\begin{array}{ccc|c} x & x & x & x \\ 0 & x & x & x \\ 0 & x & x & x \end{array} \right] \xrightarrow{i=2}$$

$$\left[\begin{array}{ccc|c} x & x & x & x \\ 0 & x & x & x \\ 0 & 0 & x & x \end{array} \right]$$

Solve by
back substitution

EXAMPLE

Solve the following system by Gaussian elimination.

$$\begin{bmatrix} 1 & 2 & -2 \\ -2 & 1 & -3 \\ 3 & 11 & 2 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ x_3 \end{bmatrix} = \begin{bmatrix} 3 \\ 6 \\ 6 \end{bmatrix}$$

$$\left[\begin{array}{ccc|c} 1 & 2 & -2 & 3 \\ -2 & 1 & -3 & 6 \\ 3 & 11 & 2 & 6 \end{array} \right] \xrightarrow{\substack{r_2 := r_2 + 2r_1 \\ r_3 := r_3 - 3r_1}} \left[\begin{array}{ccc|c} 1 & 2 & -2 & 3 \\ 0 & 5 & -7 & 12 \\ 0 & 5 & 8 & -3 \end{array} \right]$$

$$\xrightarrow{r_3 := r_3 - r_2} \left[\begin{array}{ccc|c} 1 & 2 & -2 & 3 \\ 0 & 5 & -7 & 12 \\ 0 & 0 & 15 & -15 \end{array} \right]$$

(4)

LU FACTORIZATION (WATKINS-1.7)

$$A = LU$$

$n \times n$ (nonsingular) $n \times n$ (lower triangular) $n \times n$ (upper triangular)

LU factorization can be computed by row-reducing A into U and storing multipliers

$$A = \begin{bmatrix} 2 & 3 & 1 \\ -2 & 1 & 3 \\ 6 & 1 & -1 \end{bmatrix} \xrightarrow{\substack{r_2 := r_2 + r_1 \\ r_3 := r_3 - 3r_1}} \begin{bmatrix} 2 & 3 & 1 \\ -1 & 4 & 4 \\ 3 & -8 & -4 \end{bmatrix}$$

$$\xrightarrow{r_3 := r_3 + 2r_2} \begin{bmatrix} 2 & 3 & 1 \\ -1 & 4 & 4 \\ 3 & -2 & 4 \end{bmatrix}$$

$$\underbrace{\begin{bmatrix} 1 & 0 & 0 \\ -1 & 1 & 0 \\ 3 & -2 & 1 \end{bmatrix}}_L \underbrace{\begin{bmatrix} 2 & 3 & 1 \\ 0 & 4 & 4 \\ 0 & 0 & 4 \end{bmatrix}}_U = \begin{bmatrix} 2 & 3 & 1 \\ -2 & 1 & 3 \\ 6 & 1 & -1 \end{bmatrix}$$

(unit lower triangular, i.e. entries on the diagonal are one)

FORMAL DESCRIPTION

NOTATION:

$A^{(k)}$: Matrix obtained from A
after processing columns $i=1, \dots, k-1$

$\bar{a}_j^{(k)}$: j th row of $A^{(k)}$

\bar{a}_j : j th row of A

\bar{u}_j : j th row of U

$$\begin{bmatrix} 2 & 3 & 1 \\ -2 & 1 & 3 \\ 6 & 1 & -1 \end{bmatrix} \rightsquigarrow \begin{bmatrix} 2 & 3 & 1 \\ 0 & 4 & 4 \\ 0 & -8 & -4 \end{bmatrix} \rightsquigarrow \begin{bmatrix} 2 & 3 & 1 \\ 0 & 4 & 4 \\ 0 & 0 & 4 \end{bmatrix}$$

$A = A^{(1)} \qquad A^{(2)} \qquad A^{(3)} = U$

$$\bar{a}_3^{(2)} = [0 \quad -8 \quad -4] \qquad \bar{a}_1^{(1)} = [2 \quad 3 \quad 1]$$

Note that

$$\bar{a}_1^{(1)} = \bar{u}_1 = [2 \quad 3 \quad 1]$$

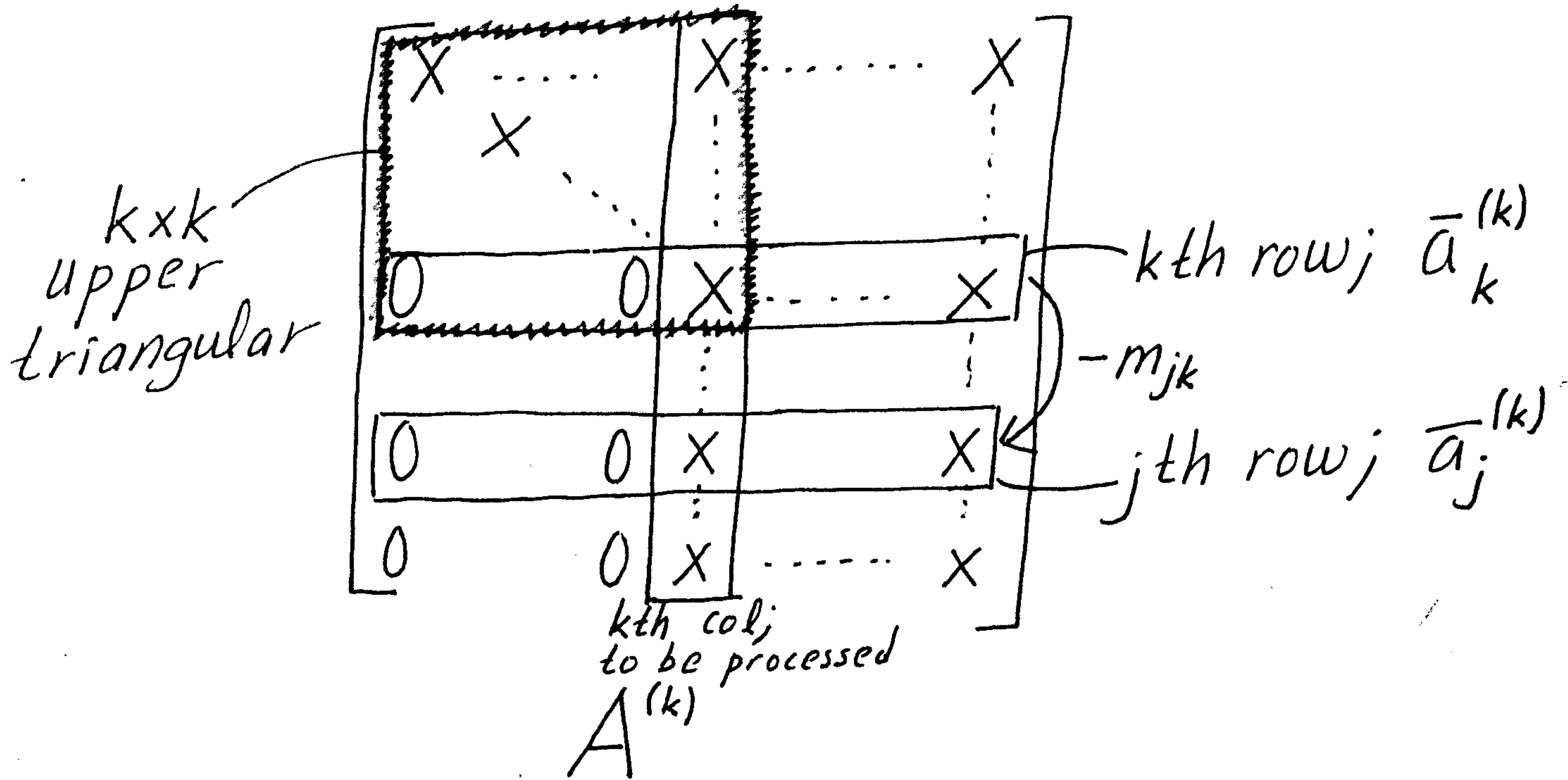
$$\bar{a}_2^{(2)} = \bar{u}_2 = [0 \quad 4 \quad 4]$$

$$\bar{a}_3^{(3)} = \bar{u}_3 = [0 \quad 0 \quad 4]$$

BASIC PROPERTIES

(i) $\bar{u}_j = a_j^{(j)} \quad j=1, \dots, n$

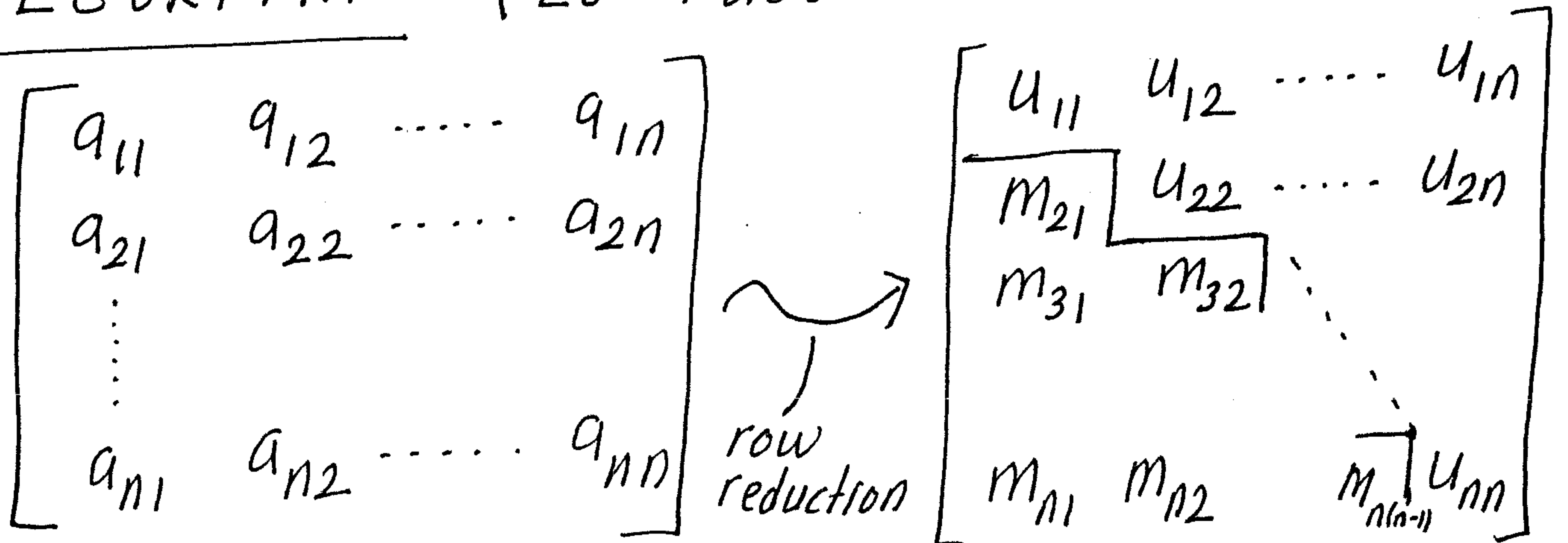
(ii) $\bar{a}_j = a_j^{(1)} \quad j=1, \dots, n$ (since $A = A^{(1)}$)



Define the multiplier m_{jk} (for $j > k$) such that

(*) $\bar{a}_j^{(k+1)} = \bar{a}_j^{(k)} - m_{jk} \bar{a}_k^{(k)}$

ALGORITHM (LU Factorization)



ASSUMPTION: Row reduction can be performed only by applying row replace operation repeatedly

PROPOSITION

$$A = \underbrace{\begin{bmatrix} 1 & & & & 0 \\ m_{21} & 1 & & & \\ m_{31} & m_{32} & \ddots & & \\ \vdots & \vdots & \ddots & \ddots & \\ m_{n1} & m_{n2} & \dots & \dots & 1 \end{bmatrix}}_L \begin{bmatrix} u_{11} & u_{12} & \dots & u_{1n} \\ & u_{22} & & u_{2n} \\ & & \ddots & \vdots \\ 0 & & & u_{nn} \end{bmatrix}$$

L (unit lower triangular; lower triangular with 1s on the diagonal)

U (upper triangular)

LEMMA

For all k, j such that $k \leq j$

$$\bar{a}_j^{(k)} = \bar{a}_j - \sum_{i=1}^{k-1} m_{ji} \bar{u}_i.$$

PROOF

The proof is by induction on k . (j is fixed)

Base Case ($k=1$)

$$\bar{a}_j^{(1)} = \bar{a}_j \quad \left(\begin{array}{l} \text{By property (ii),} \\ \text{since } A = A^{(1)} \end{array} \right)$$

Inductive Case ($k \geq 2$)

$$\text{Assume } \bar{a}_j^{(k-1)} = \bar{a}_j - \sum_{i=1}^{k-2} m_{ji} \bar{u}_i \quad \left(\begin{array}{l} \text{Inductive} \\ \text{hypothesis} \end{array} \right)$$

By equation (*) on page (7)

$$\bar{a}_j^{(k)} = \bar{a}_j^{(k-1)} - m_{j(k-1)} \bar{a}_{k-1}^{(k-1)}$$

(8)

From the inductive hypothesis and property (i)

$$\begin{aligned}
 a_j^{(k)} &= \underbrace{\left(\bar{a}_j - \sum_{i=1}^{k-2} m_{ji} \bar{u}_i \right)}_{a_j^{(k-1)}} - \underbrace{m_{j(k-1)} \bar{u}_{k-1}}_{m_{j(k-1)} \bar{a}_{k-1}^{(k-1)}} \\
 &= \bar{a}_j - \sum_{i=1}^{k-1} m_{ji} \bar{u}_i
 \end{aligned}$$

□

PROOF OF THE PROPOSITION

Apply lemma for $k=j$ to obtain

$$\frac{\bar{a}_j^{(j)}}{\bar{u}_j} = \bar{a}_j - \sum_{i=1}^{j-1} m_{ji} \bar{u}_i$$

$$\implies \bar{a}_j = \sum_{i=1}^{j-1} m_{ji} \bar{u}_i + \bar{u}_j$$

$$= \begin{bmatrix} m_{j1} & m_{j2} & \cdots & m_{j(j-1)} & 1 \end{bmatrix} \begin{bmatrix} \bar{u}_1 \\ \bar{u}_2 \\ \vdots \\ \bar{u}_j \end{bmatrix}$$

$$= \begin{bmatrix} m_{j1} & m_{j2} & \cdots & m_{j(j-1)} & 1 & 0 & \cdots & 0 \end{bmatrix} \begin{bmatrix} \bar{u}_1 \\ \bar{u}_2 \\ \vdots \\ \bar{u}_n \end{bmatrix}$$

$\underbrace{\hspace{10em}}_U \textcircled{9}$

Combine the rows of A

$$A = \begin{bmatrix} \bar{a}_1 \\ \bar{a}_2 \\ \vdots \\ \bar{a}_n \end{bmatrix} = \begin{bmatrix} 1 & & & & \\ m_{21} & 1 & & & 0 \\ \vdots & \vdots & \ddots & & \vdots \\ m_{n1} & m_{n2} & \cdots & m_{n(n-1)} & 1 \end{bmatrix} \begin{bmatrix} \bar{u}_1 \\ \bar{u}_2 \\ \vdots \\ \bar{u}_n \end{bmatrix}$$

$\underbrace{\hspace{15em}}_L \qquad \underbrace{\hspace{2em}}_U$



PSEUDOCODE TO COMPUTE LU FACTORIZATION

NOTATION

$a_{j, m_1:m_2}$: j th row of A restricted to the entries with column numbers in $[m_1, m_2]$

e.g. $A = \begin{bmatrix} 5 & -2 & 7 & 8 \\ 1 & -3 & 2 & 4 \end{bmatrix}$ $a_{2, 2:3} = [-3 \ 2]$
 $a_{2, 2:4} = [-3 \ 2 \ 4]$

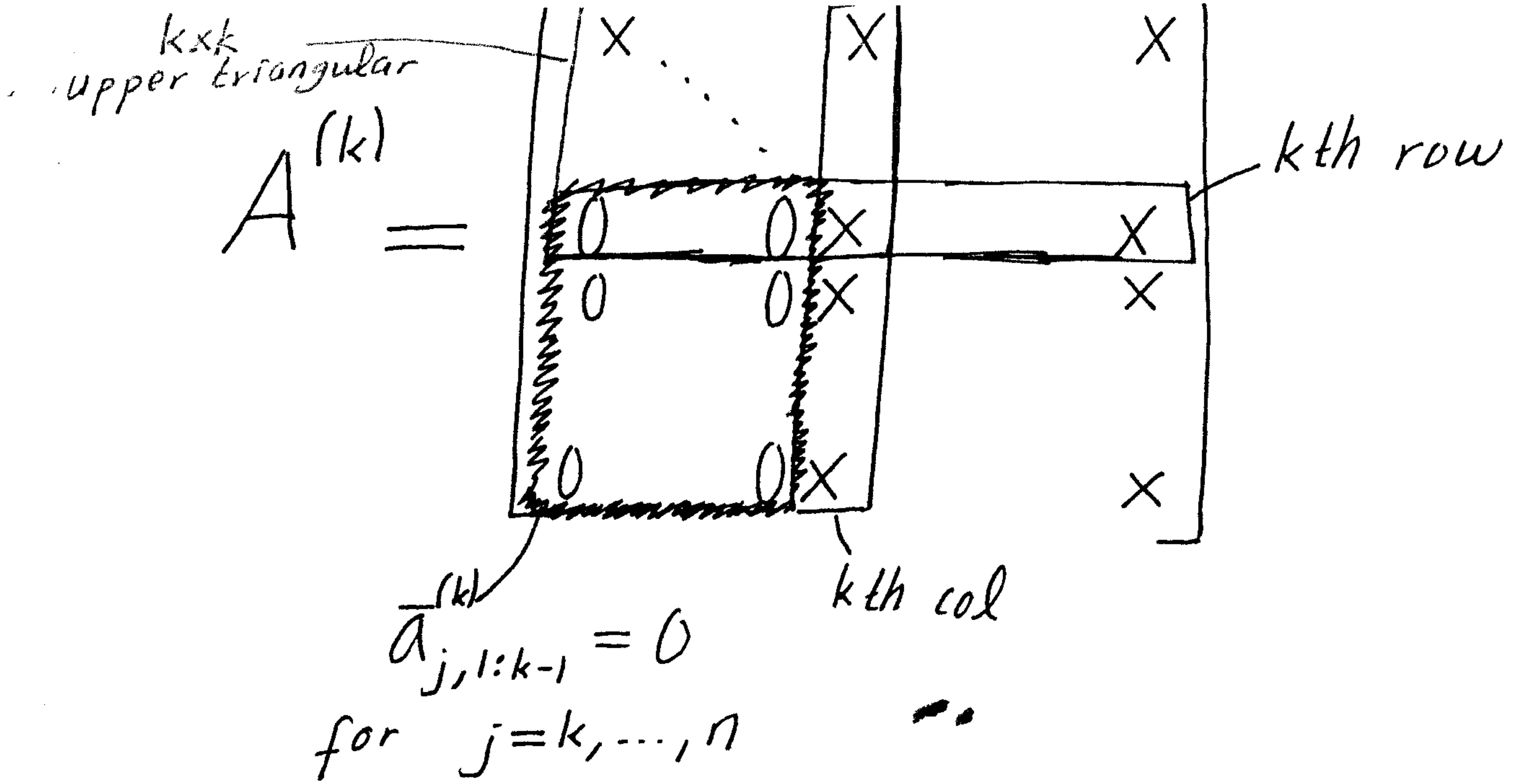
To process k th column we need to perform

$$(**) \ a_j^{(k+1)} = a_j^{(k)} - m_{jk} a_k^{(k)} \quad \text{for } j = k+1, \dots, n$$

But note

$$(1) \ a_{j, 1:k-1}^{(k)} = 0 \quad \text{for } j = k, k+1, \dots, n$$

$$(2) \ \text{we choose } m_{jk} \text{ so that } a_{j,k}^{(k+1)} = 0 \quad \text{for } j = k+1, \dots, n$$



Instead of $(**)$ perform

$$a_{j, k+1:n}^{(k+1)} = a_{j, k+1:n}^{(k)} - m_{jk} a_{k, k+1:n}^{(k)}$$

PSEUDOCODE (LU Factorization)

for $k = 1, \dots, n-1$ (k : col to be processed)

for $j = k+1, \dots, n$ (j : row number)

$m_{jk} \leftarrow a_{jk} / a_{kk}$ } 1 flop

$a_{j, k+1:n} \leftarrow a_{j, k+1:n} - m_{jk} a_{k, k+1:n}$ } $2(n-k)$ flops

$a_{jk} \leftarrow m_{jk}$ (store multiplier below the main diagonal)

end

end

OPERATION COUNT

For $k=1, \dots, n-1$ and $j=k+1, \dots, n$

(1) a division is required for m_{jk}

(2) $2(n-k)$ flops are required to perform

$$a_{j,k+1:n} \leftarrow a_{j,k+1:n} - m_{jk} a_{k,k+1:n}$$

$$\text{TOTAL FLOPS} = \sum_{k=1}^{n-1} \sum_{j=k+1}^n (2(n-k) + 1)$$

$$= \sum_{k=1}^{n-1} (2(n-k)^2 + (n-k))$$

$$= 2 \left((n-1)^2 + (n-2)^2 + \dots + 1^2 \right)$$

+

$$\left((n-1) + (n-2) + \dots + 1 \right)$$

$$= 2 \frac{(n-1)n(2n-1)}{6} + \frac{n(n-1)}{2}$$

$$= \frac{2n^3}{3} + O(n^2)$$

LU FACTORIZATION TO SOLVE LINEAR SYSTEMS

Consider

$$\begin{bmatrix} 2 & 3 & 1 \\ -2 & 1 & 3 \\ 6 & 1 & -1 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ x_3 \end{bmatrix} = \begin{bmatrix} 9 \\ -5 \\ 11 \end{bmatrix}$$

equivalently

$$\begin{bmatrix} 1 & 0 & 0 \\ -1 & 1 & 0 \\ 3 & -2 & 1 \end{bmatrix} \underbrace{\begin{bmatrix} 2 & 3 & 1 \\ 0 & 4 & 4 \\ 0 & 0 & 4 \end{bmatrix}}_{\hat{X}} \begin{bmatrix} x_1 \\ x_2 \\ x_3 \end{bmatrix} = \begin{bmatrix} 9 \\ -5 \\ 11 \end{bmatrix}$$

(1) FORWARD STAGE

$$\begin{bmatrix} 1 & 0 & 0 \\ -1 & 1 & 0 \\ 3 & -2 & 1 \end{bmatrix} \begin{bmatrix} \hat{x}_1 \\ \hat{x}_2 \\ \hat{x}_3 \end{bmatrix} = \begin{bmatrix} 9 \\ -5 \\ 11 \end{bmatrix}$$

Solve by forward substitution

$$\hat{x}_1 = 9, \hat{x}_2 = 4, \hat{x}_3 = -8$$

(2) BACKWARD STAGE

$$\begin{bmatrix} 2 & 3 & 1 \\ 0 & 4 & 4 \\ 0 & 0 & 4 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ x_3 \end{bmatrix} = \begin{bmatrix} 9 \\ 4 \\ -8 \end{bmatrix}$$

Solve by back substitution

$$x_3 = -2, x_2 = 3, x_1 = 1$$

PROCEDURE (To solve $Ax=b$)

FLOPS

(1) Compute the LU-factorization

$$A = LU$$

$$\frac{2n^3}{3} + O(n^2)$$

(2) Define $\hat{x} := Ux$. Solve

by $L\hat{x} = b$
forward substitution

$$n^2$$

(3) Solve

by $Ux = \hat{x}$
back substitution

$$n^2$$