

# LECTURE 25

## QR ALGORITHM WITH SHIFTS

QR algorithm, when it converges, converges to an upper triangular matrix only  $q$ -linearly, i.e.

$$\overbrace{Q_k^* \dots Q_1^* \underbrace{A}_{\text{Hessenberg}} Q_1 \dots Q_k}^{A_k} \longrightarrow \underbrace{T}_{\text{upper triangular}} \text{ as } k \rightarrow \infty$$

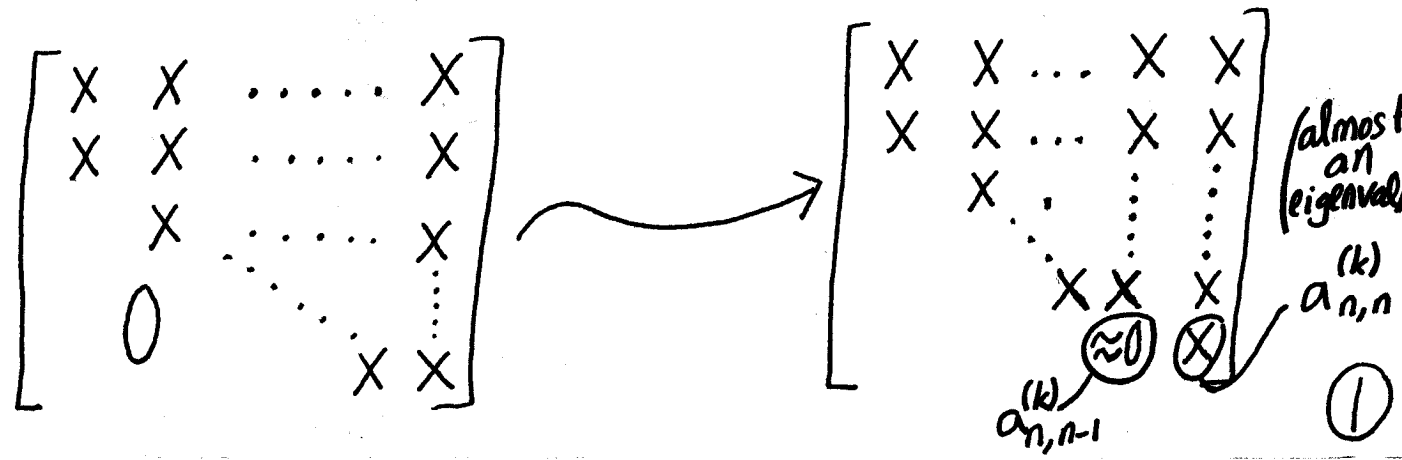
and subdiagonal entries of  $A$  approach zero  $q$ -linearly. For instance typically

$$\lim_{k \rightarrow \infty} \frac{|a_{n,n-1}^{(k+1)} - 0|}{|a_{n,n-1}^{(k)} - 0|} = \frac{|\lambda_n|}{|\lambda_{n-1}|}$$

$a_{n,n-1}^{(k)}$ :  $(n, n-1)$  entry of  $A_k$

$\lambda_n$ : Smallest eigenvalue of  $A$  in modulus

$\lambda_{n-1}$ : Second smallest eigenvalue of  $A$  in modulus



To speed up convergence shift  $A_k$  by  $M_k = a_{n,n}^{(k)}$

Define the sequence  $\{A_k\}$  such that

$$(i) A_0 = A$$

For  $k \geq 0$  (ii)  $(A_k - M_k I) = Q_{k+1} R_{k+1}$  and  $A_{k+1} = R_{k+1} Q_{k+1} + M_k I$

### ALGORITHM (QR Algorithm with Shifts)

\* Given  $A \in \mathbb{C}^{n \times n}$  in Hessenberg form.

\* Produce a sequence  $\{A_k\}$  such that typically  $\lim_{k \rightarrow \infty} A_k$  is upper triangular.

$$A_0 = A$$

for  $k = 0, 1, \dots$

Choose a shift  $M_k$

Compute a QR factorization  $A_k - M_k I = Q_{k+1} R_{k+1}$

end  $A_{k+1} = R_{k+1} Q_{k+1} + M_k I$

### COMMON SHIFT STRATEGIES

Rayleigh shift:  $M_k = a_{n,n}^{(k)}$

Wilkinson shift:  $M_k$  is the eigenvalue of  $A_k(n-1:n, n-1:n)$  closest to  $a_{n,n}^{(k)}$

## EXAMPLE

$$A = \begin{bmatrix} 3 & 4 \\ 2 & 1 \end{bmatrix} \text{ has eigenvalues } \lambda_1 = 5, \lambda_2 = -1$$

Apply QR algorithm with fixed shifts  $\mu = -0.8$ .

$$Q_1 R_1 = A + 0.8I$$

$$A_1 = R_1 Q_1 - 0.8I = \begin{bmatrix} 5.04 & -1.87 \\ 0.13 & -1.04 \end{bmatrix}$$

$$Q_2 R_2 = A_1 + 0.8I$$

$$A_2 = R_2 Q_2 - 0.8I = \begin{bmatrix} 5.00 & 2.00 \\ 0.00 & -1.00 \end{bmatrix}$$

## REMARKS

(i)  $A_{k+1}$  and  $A_k$  are unitarily similar.

$$(A_k - \mu_k I) = Q_{k+1} R_{k+1} \text{ and } A_{k+1} = R_{k+1} Q_{k+1} + \mu_k I$$

$$\implies R_{k+1} = Q_{k+1}^* (A_k - \mu_k I) \text{ and } A_{k+1} = R_{k+1} Q_{k+1} + \mu_k I$$

$$\begin{aligned} A_{k+1} &= Q_{k+1}^* (A_k - \mu_k I) Q_{k+1} + \mu_k I \\ &= Q_{k+1}^* A_k Q_{k+1} \end{aligned}$$

(ii) Rate of convergence with Rayleigh or Wilkinson shifts is typically  $q$ -quadratic. (3)

## QR ALGORITHM WITH DOUBLE SHIFTS

In practice Wilkinson shift is the standard shift strategy used.

When the matrix  $A$  is real, ideally we like to remain in real arithmetic, i.e. ideally we like  $A_k$  to be real for all  $k$ .

The Wilkinson shift  $\mu_k$  can be complex, but the matrix

$$(A_k - \mu_k I)(A_k - \bar{\mu}_k I)$$

is real provided  $A_k$  is real.

$$A_k^2 - (\mu_k + \bar{\mu}_k)A_k + |\mu_k|^2 I$$

The standard double-shift strategy with Wilkinson shifts  $\mu_k$  generates a sequence  $\{A_k\}$  such that

(i)  $A_0 = A$

(ii)  $(A_k - \mu_k I) = Q_{k+1} R_{k+1}$  and  $A_{k+1} = R_{k+1} Q_{k+1} + \mu_k I$

(iii)  $(A_{k+1} - \bar{\mu}_k I) = Q_{k+2} R_{k+2}$  and  $A_{k+2} = R_{k+2} Q_{k+2} + \bar{\mu}_k I$

FOR EVEN  $k \geq 0$

(4)

# REMARK

$$A_{k+2} = Q_{k+2}^* Q_{k+1}^* A_k Q_{k+1} Q_{k+2} \text{ is real.}$$

since

$$(A_k - \bar{\mu}_k I) (A_k - \mu_k I)$$

$$\stackrel{=}{=} (A_k - \bar{\mu}_k I) Q_{k+1} R_{k+1}$$

$$\stackrel{=}{=} Q_{k+1} (A_{k+1} - \bar{\mu}_k I) R_{k+1}$$

$$\stackrel{=}{=} Q_{k+1} Q_{k+2} R_{k+2} R_{k+1}$$

QR factorization

of  $(A_k - \bar{\mu}_k I)(A_k - \mu_k I)$

implying  $Q_{k+1} Q_{k+2}$  is real.

NOTE

$$A_{k+1} = Q_{k+1}^* A_k Q_{k+1}$$

$$\stackrel{\Rightarrow}{=} (A_{k+1} - \bar{\mu}_k I)$$

$$Q_{k+1}^* (A_k - \bar{\mu}_k I) Q_{k+1}$$

$$\stackrel{\Rightarrow}{=} Q_{k+1} (A_{k+1} - \bar{\mu}_k I) Q_{k+1}$$

# SIMULTANEOUS POWER ITERATION

## Power Iteration

$$z_k = A^k q_0$$

$$q_k = z_k / \|z_k\|$$

$$\hat{\lambda}_k = q_k^* A q_k$$

## Simultaneous Power Iteration

$$Z_k = A^k Q_0$$

$\underbrace{\quad}_{n \times n} \quad \underbrace{\quad}_{n \times n}$

$$Q_k \text{ satisfies } Z_k = \underbrace{Q_k R_k}_{\text{QR factor}}$$

$$\hat{\Lambda}_k = Q_k^* A Q_k$$

For simplicity assume  $A \in \mathbb{C}^{n \times n}$  is Hermitian.

## Power Iteration

(1)  $q_k$  converges to a dominant eigenvector

(2)  $\hat{\lambda}_k$  converges to the eigenvalue with largest modulus

## Simultaneous Power Iteration

WILL BE DISCUSSED

$Q_k \rightarrow V = [v_1 \dots v_n]$  as  $k \rightarrow \infty$  where  $v_1, \dots, v_n$  are orthonormal full set of eigenvectors

$\hat{\Lambda}_k$  converges to the full set of eigenvalues

As will be shown the QR algorithm is equivalent to simultaneous power iteration. (with  $Q_0 = I$ )

In particular  $A_k = \hat{\Lambda}_k (= Q_k^* A Q_k)$

estimates by QR algorithm

# EXAMPLE

$$A = \begin{bmatrix} 3 & 4 \\ 2 & 1 \end{bmatrix} \text{ with eigenvalues } \lambda_1 = 5, \lambda_2 = -1$$

Recall when the QR algorithm is applied

$$A_3 = R_3 Q_3 = \begin{bmatrix} 5.01 & -1.97 \\ 0.03 & -1.01 \end{bmatrix}$$

Apply the simultaneous power iteration  
with  $Q_0 = I$

$$* Z_3 = A^3(I) = \begin{bmatrix} 83 & 84 \\ 42 & 41 \end{bmatrix}$$

$$* Q_3 = \begin{bmatrix} -0.89 & -0.45 \\ -0.45 & 0.89 \end{bmatrix} \text{ satisfies}$$

$$Z_3 = \underbrace{Q_3 R_3}_{\text{QR Factorization}}$$

$$* \Lambda_3 = Q_3^* A Q_3 = \begin{bmatrix} 5.01 & -1.97 \\ 0.03 & -1.01 \end{bmatrix}$$

# SUBSPACE ITERATION

Denote

\* the eigenvalues of  $A$  by  
 $\lambda_1, \dots, \lambda_n$  s.t.

$$|\lambda_1| \geq |\lambda_2| \geq \dots \geq |\lambda_n|,$$

\* the associated ~~eigenvalues~~ <sup>(orthonormal)</sup> eigenvectors  
by  $v_1, v_2, \dots, v_n$ .

Given  $q_1, \dots, q_m \in \mathbb{C}^n$  with  $m \leq n$ . Let  
 $S = \text{span} \{q_1, \dots, q_m\}$ ,

and

$$AS = \{Ax : x \in S\} \quad \left( \begin{array}{l} \text{the image of} \\ S \text{ under the} \\ \text{transformation} \\ x \rightarrow Ax \end{array} \right)$$

## REMARK

$AS$  is a subspace of  $\mathbb{R}^n$ .

## THM (Subspace Iteration)

Let  $U_m = \text{span} \{v_{m+1}, \dots, v_n\}$  and

$S_k = A^k S$ . If  $S \cap U_m = \{0\}$  and  $|\lambda_m| > |\lambda_{m+1}|$ ,  
then

$$S_k \longrightarrow T_m \quad \text{as } k \rightarrow \infty$$

where  $T_m = \text{span} \{v_1, \dots, v_m\}$ .



# CONVERGENCE OF SIMULTANEOUS ITERATION

Letting  $S^{(m)} = \text{span}\{e_1, \dots, e_m\} \subseteq \mathbb{C}^n$   
and denoting the image of  $S^{(m)}$  under  
the transformation  $x \rightarrow A^k x$  by

$$S_k^{(m)} = \{A^k x : x \in S^{(m)}\}$$

we deduce (by thm-subspace iteration)

$$S_k^{(m)} \rightarrow T_m \quad \text{as } k \rightarrow \infty.$$

(Assuming  $A$  has distinct eigenvalues  
and  $S^{(m)} \cap U_m = \{0\}$ .)

Consequently

$$\underbrace{\text{Range}(A^k(:, 1:m))}_{S_k^{(m)}} \rightarrow \underbrace{\text{span}\{v_1, \dots, v_m\}}_{T_m} \quad \text{as } k \rightarrow \infty$$

QR factorization of  $A^k(:, 1:m)$

$$\begin{aligned} A^k(:, 1:m) &= QR \\ &= \begin{matrix} n \times m & m \times m \\ \underbrace{[q_1^{(k)} \dots q_m^{(k)}]} & R \end{matrix} \end{aligned}$$

## Note

Recall that  $\{q_1^{(k)}, \dots, q_l^{(k)}\}$  is an orthonormal basis for  $\text{range}(A^k(:, 1:l))$  for  $l=1, \dots, m$

(i)  $\text{Range}(A^k(:, 1)) \longrightarrow \text{span}\{v_1\}$  as  $k \rightarrow \infty$ .

Furthermore

$\{q_1^{(k)}\}$  is a basis for  $\text{Range}(A^k(:, 1))$

$$q_1^{(k)} \xrightarrow{\implies} \underbrace{c_1}_{\text{scalar}} v_1 \text{ as } k \rightarrow \infty$$

(ii)  $\text{Range}(A^k(:, 1:2)) \longrightarrow \text{span}\{v_1, v_2\}$  as  $k \rightarrow \infty$ .

Furthermore

$\{q_1^{(k)}, q_2^{(k)}\}$  is an orthonormal basis for  $\text{Range}(A^k(:, 1:2))$

$$\xrightarrow{\implies} (q_2^{(k)} \perp q_1^{(k)} \text{ and } q_1^{(k)} = c_1 v_1)$$
$$q_2^{(k)} \xrightarrow{\implies} \underbrace{c_2}_{\text{scalar}} v_2 \text{ as } k \rightarrow \infty$$

(iii)  $\text{Range}(A^k(:, 1:l)) \longrightarrow \text{span}\{v_1, \dots, v_l\}$  as  $k \rightarrow \infty$  for  $l=1, \dots, m$ .

$\{q_1^{(k)}, \dots, q_l^{(k)}\}$  is an orthonormal basis for  $\text{Range}(A^k(:, 1:l))$

$$\xrightarrow{\implies} (q_l^{(k)} \perp q_j^{(k)} \text{ and } q_j^{(k)} = c_j v_j \text{ for } j < l)$$
$$q_l^{(k)} \xrightarrow{\implies} c_l v_l \text{ as } k \rightarrow \infty$$

## THM (Simultaneous Iteration)

Let  $A \in \mathbb{C}^{n \times n}$  be a Hermitian matrix with distinct eigenvalues and

$$A^k = Q_k R_k$$

be a QR factorization. Then generically the columns of  $Q_k$  approach the eigenvectors  $v_1, v_2, \dots, v_n$  of  $A$  as  $k \rightarrow \infty$ .